

APPARATUS AND METHOD FOR VERIFYING MEMORY COHERENCY OF DUPLICATION PROCESSOR

BACKGROUND OF THE INVENTION

1. Field of the Invention

The present invention relates to a duplication processor of a switching system, and more particularly to an apparatus and a method for verifying memory coherency of a duplication processor.

2. Description of the Background Art

Generally, a processor of a switching system has a duplication structure of two processors, having the same construction to each other, that is, an active processor which is operated in an active mode and a standby processor which is operated in a standby mode

Figure 1 is a schematic block diagram of a duplication processor including an active processor 10 and a standby processor 20.

As shown in the drawing, the active processor 10 includes a first CPU 11, a first memory controller 12, a local bus controller 13 and a first duplication pathway 14.

The first CPU 11 performs a general controlling operation of the processor, and for verifying memory coherency, it starts a standby memory read transaction (SMRT) and informs the local bus controller 13 of it.

The first memory controller 12 arbitrates a processor bus (P-bus) according to a transaction request by the local bus controller 13.

active processor starts reading operation of the standby memory, a long time is taken for transaction according to an operation state of the D-ch and the standby processor. As a result, the CPU of the active processor should wait to proceed to other operation until the transaction operation is completely performed.

5 In addition, in the conventional duplication processor structure, each transaction is performed by a single beat, resulting in that a long time is to be taken to perform coherency verifying for all the standby memories,

SUMMARY OF THE INVENTION

10 Therefore, an object of the present invention is to provide an apparatus and method for verifying memory coherency of duplication processor which is capable of improving a use efficiency of a CPU by reducing a load at an active processor generated when a standby memory is read.

15 Another object of the present invention is to provide an apparatus and method for verifying memory coherency of duplication processor which is capable of reducing the time of coherency verification by providing an active processor with a cacheability of a standby memory image.

To achieve these and other advantages and in accordance with the
20 purpose of the present invention, as embodied and broadly described herein, there is provided an apparatus for verifying memory coherency of a duplication processor having a symmetrical structure including: an active processor in which a standby memory read command (SMRC) is generated and transmitted by hardware and then a read data of the standby memory which has been inputted
25 corresponding to the SMRC is image-buffered to verify a memory coherency; and

a standby processor in which the SMRC transmitted from the active processor is analyzed and a read command of a standby memory is outputted, and then the data read from the standby memory is transmitted to the active processor.

In order to achieve the above objects, the active processor of the apparatus for verifying memory coherency of a duplication processor includes: a standby memory image buffer (SMIB) for temporarily storing a read data of a standby memory; a CPU for generating an SMRC; a first memory controller for storing the read data of the standby memory in the SMIB; a first processor bus controller for applying the read data of the standby memory to the first memory controller; and a first duplication processor for informing the standby processor of a registered SMRC when the SMRC is registered by the CPU, and outputting a command done signal to the CPU when the reading operation of the standby memory is completed.

In order to achieve the above objects, the standby processor of the apparatus for verifying memory coherency of a duplication processor includes: a second memory controller for controlling access of the standby memory; a second processor bus controller for transmitting a received SMRC to the second memory controller, and a second duplication processor for analyzing the SMRC inputted through the second memory controller, sequentially generating a read address of the standby memory, and transmitting the data read from the standby memory to the active processor.

In the present invention, the SMRC includes a start address of the standby memory, a size of a data to be read and an address of the SMRC to store the read data.

In order to achieve the above objects, there is also provided a method for

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accompanying drawings.

BRIEF DESCRIPTION OF THE DRAWINGS

The accompanying drawings, which are included to provide a further
5 understanding of the invention and are incorporated in and constitute a part of this
specification, illustrate embodiments of the invention and together with the
description serve to explain the principles of the invention.

In the drawings:

Figure 1 is a schematic block diagram showing a construction of a
10 duplication processor in accordance with a conventional art;

Figure 2 is a schematic block diagram showing a construction of a
duplication processor in accordance with the present invention;

Figure 3 is a flow chart of a method for verifying memory coherency of the
duplication processor in accordance with the present invention;

15 Figure 4 is a flow chart of an operation of a standby processor of Figure 3
in accordance with the present invention; and . . .

Figure 5 is a flow chart of an operation of a standby processor of Figure 3
in accordance with the present invention.

20 DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Reference will now be made in detail to the preferred embodiments of the
present invention, examples of which are illustrated in the accompanying drawings

The present invention is featured in that a reading operation of a standby
25 memory, which has been performed by software in the conventional art, is

performed by hardware, and a burst mechanism is provided for an image of the standby memory. That is, a CPU of an active processor of the present invention registers only a standby memory read command (SMRC) in the duplication processor, and performs a different operation until the memory reading operation is completed.

In addition, the present invention is also featured in that a standby memory image buffer (SMIB) different to an actual standby memory is implemented in a predetermined region of an active memory, so that a burst transaction can be performed in verifying memory coherency through the reading operation of the standby memory.

Figure 2 is a schematic block diagram showing a construction of a duplication processor in accordance with the present invention.

As shown in the drawing, the duplication processor of the present invention additionally includes a first and a second duplication processor 140 and 230, and a standby memory image buffer (SMIB) 130 into the construction of the conventional art as shown in Figure 1.

A CPU 110 of an active processor 100 controls the overall system, and generates a standby memory read command (SMRC) in verifying coherency for the standby memory. At this time, the SMRC includes a start address of a standby memory 220, the data size and an address indicating an SMIB in which a read data is to be stored.

A first memory controller 120 gives a bus use right to a first processor bus controller 150 and stores a read data of the standby memory 220 in the SMIB 130.

The SMIB (Standby Memory Image Buffer) 130 is a block for temporarily storing the read data of the standby memory 220. The present invention includes

two SMIBs 130, of which one SMIB is used to verify coherency, the other is used to store the read data of the standby memory 220. Each SMIB is discriminated by 1bit SMIB address included in the SMRC.

5 A first duplication processor 140 transmits the SMRC outputted from the CPU 110 to the standby processor, and outputs a command done signal (CDS) to the CPU 110 according to a write done signal (WDS) outputted from the first processor bus controller 150.

10 The first processor bus controller 150 transmits the read data of the standby memory 220 to the first memory controller 120, and when the final read data is received, the first processor bus controller 150 outputs a write done signal (WDS) to the first duplication processor 140.

The first duplication pathway 160 and the active memory 170 are performed the same manner as in the conventional art, descriptions of which are thus omitted.

15 A second memory controller 210 of a standby processor 200 gives a bus use right to a second processor bus controller 240, and transmits the SMRC transmitted from the active processor 100 to a second duplication processor 230.

A second memory controller 210 accesses to the standby memory 220 according to a read address (RA) outputted from the second duplication processor 20 230.

The second duplication processor 230 sequentially generates the read address (RA) of the standby memory 220 according to the start address and the data size included in the SMRC, and attaches an SMIB address to each read data of the standby memory 220.

25 A standby memory 220, a second processor bus controller 240 and a

second duplication pathway 250 of the present invention are the same as that of the conventional art, so that descriptions therefor are omitted.

The operation of the duplication processor constructed as described above will now be explained.

5 Figure 3 is a flow chart of a method for verifying memory coherency of the duplication processor in accordance with the present invention.

As shown in the drawing, in the present invention, the SMRC registering process (S1), the SMRC processing process (S2) and the coherency verifying process (S3) are sequentially performed.

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1) SMRC registering:

First, the CPU 110 of the active processor 100 generates an SMRC including a start address, the size of a data to be read and an SMIB address for writing a read data, and outputs the SMRC to the first duplication processor 140, thereby registering the SMRC (S1). After the SMRC is registered, the SMRC is processed by the first duplication processor 140, so that the CPU 110 can perform other operation while the reading operation of the standby memory 220 is performed.

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2) SMRC processing

The first duplication processor 140 of the active processor 100 outputs the inputted SMRC through the first duplication pathway 160 to the standby processor 200. Then, the second duplication processor 230 of the standby processor 200 receives the SMRC through the second duplication pathway 250, the second processor bus controller 240 and the second memory controller 210, and

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generates a read address (RA) of the standby memory 220 on the basis of information included in the SMRC.

In detail, the second duplication processor 230 analyzes the SMRC (S10) and sequentially generates read addresses (RA) of the standby memory 220 on the basis of the start address and the size of the data included in the SMRC. For example, assuming that a start address is 'A' and the size of the data is '4', the second duplication processor 230 sequentially generates read addresses from 'A' to 'D' (S11). Then, the second memory controller 210 reads data corresponding to the read addresses (RA) from the standby memory 220 and outputs them to the second duplication processor 230 (S12).

Whenever the second duplication processor 230 receives the read data of the standby memory 220 from the second memory controller 210, the second duplication processor 230 attaches an SMIB address to each read data and transmits it through the second duplication pathway 250 to the active processor 100 (S13, S14). And, at this time, the second duplication processor 230 confirms whether the data has been all read as much as the size set in the SMRC (S15), and if the data as much as the set size is all read, the second duplication processor 230 sets a transmission completion flag to the final read data and transmits it through the first duplication pathway 250 to the active processor 100 (S16).

When the read data of the standby memory 220 is received through the first duplication pathway 160 (S20), the first processor controller 150 checks the SMIB address attached to each read data and stores the corresponding read data in the SMIB 130 (S21, S22). And, at this time, the first processor controller 150 checks a transmission completion flag of each read data (S23), and if the

transmission completion flag has been set, the first processor controller 150 stores the final read data in the SMIB 130 and outputs the write done signal (WDS) to the first duplication processor 140 (S24).

Then, the first duplication processor 140 generates a command done signal (CDS) according to the write done signal (WDS) outputted from the first processor bus controller 150 and outputs it to the CPU 110 (S25), thereby informing that the data of the standby memory 220 desired by the CPU 110 has been completely stored in the SMIB 130.

10 3) Coherency verification

When the CPU 110 receives the command done signal (CDS) from the first duplication processor 140, it compares the data stored in SMIB 130 and the data stored in the active memory 170 to perform a memory coherency verifying operation.

15 As so far described, according to the apparatus and method for verifying memory coherency of a duplication processor of the present invention, the CPU of the active processor outputs only the SMRC command and the SMRC is separately processed by the duplication processor. Thus, the load of the CPU is reduced, and accordingly, the use efficiency of the CPU can be increased.

20 In addition, the burst transaction can be performed both when the data is read from the standby memory and when the read data is transmitted, so that the use efficiency of the processor bus, the duplication bus and the duplication channel can be improved. Especially, a bad influence according to the operation of each processor and the duplication channel can be minimized.

25 Moreover, the SMIB provides burst mechanism function, so that the CPU

of the active processor can perform the burst transaction in verifying a memory coherency. Thus, the time required for verification can be remarkably reduced.

As the present invention may be embodied in several forms without departing from the spirit or essential characteristics thereof, it should also be understood that the above-described embodiments are not limited by any of the details of the foregoing description, unless otherwise specified, but rather should be construed broadly within its spirit and scope as defined in the appended claims, and therefore all changes and modifications that fall within the meets and bounds of the claims, or equivalence of such meets and bounds are therefore intended to be embraced by the appended claims.